

CPSC 213 Lab 3

Dynamic Arrays & C Pointers

Slides available at randyzhu.com/cpsc213

Course Updates

Quiz 1; A3 due Friday

- Assignment 3 is due on Friday, January 30th
- Quiz 1 is running today and tomorrow
 - No questions *about* quizzes

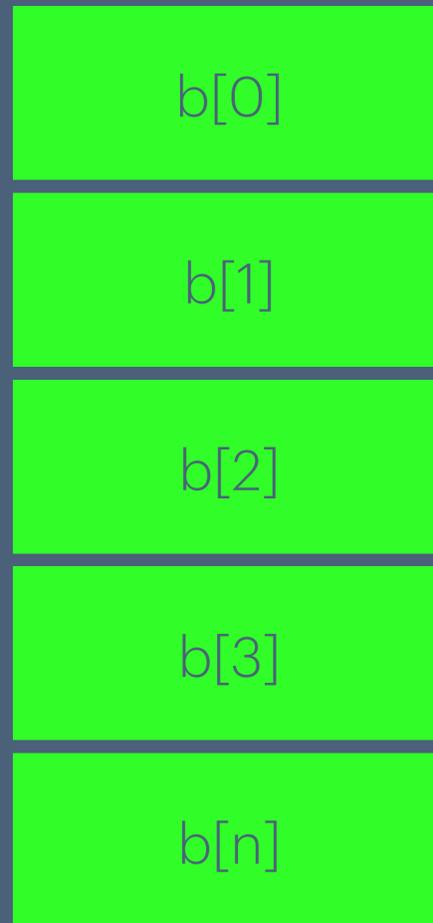
Pointer Review

What is a pointer?

- A variable that contains a memory address
- Syntax examples:
 - `int* ptr;`
 - `unsigned char* m;`
 - `int** matrix;`
- in general: `<type>* ptr;`
- Pointers can point to any type, *including pointers*

```
int* b;  
int a[10];
```

0x????:

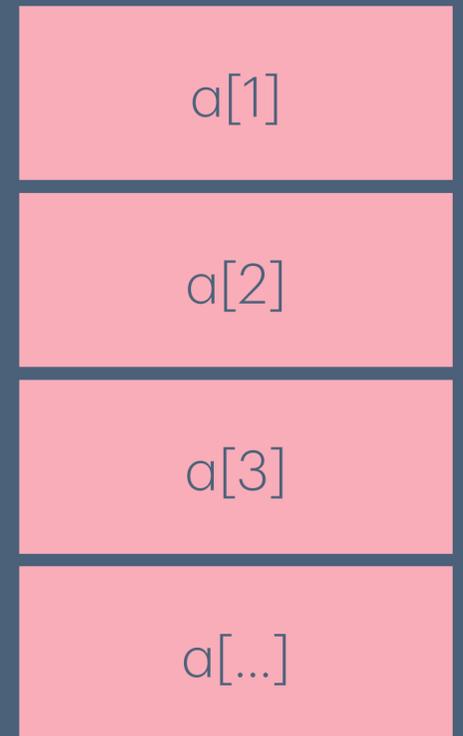


0x1000:

b:



a:



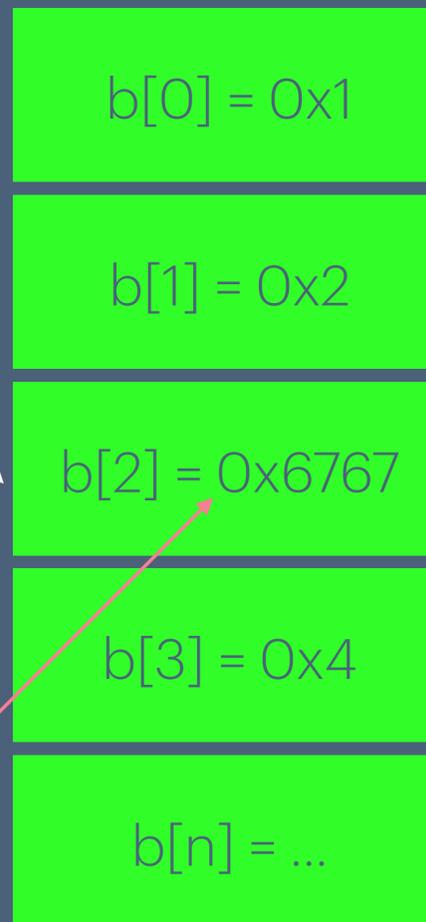
Pointer Review

Array-like; pointer arithmetic

- `*b` gets the value a pointer points to
- Here, `*b` \rightarrow `0x1`
- `b + 1`
 - Gets the address of "next" int after b
 - $0x4000 + \text{sizeof}(\text{int}) * 1 = 0x4000 + 4 = \mathbf{0x4004}$
- `*(b + 1)` gets the value of next int after b
 - Equivalent to `b[1]`
 - In general `*(b + n) = b[n]`

```
int* b;  
int a[10];
```

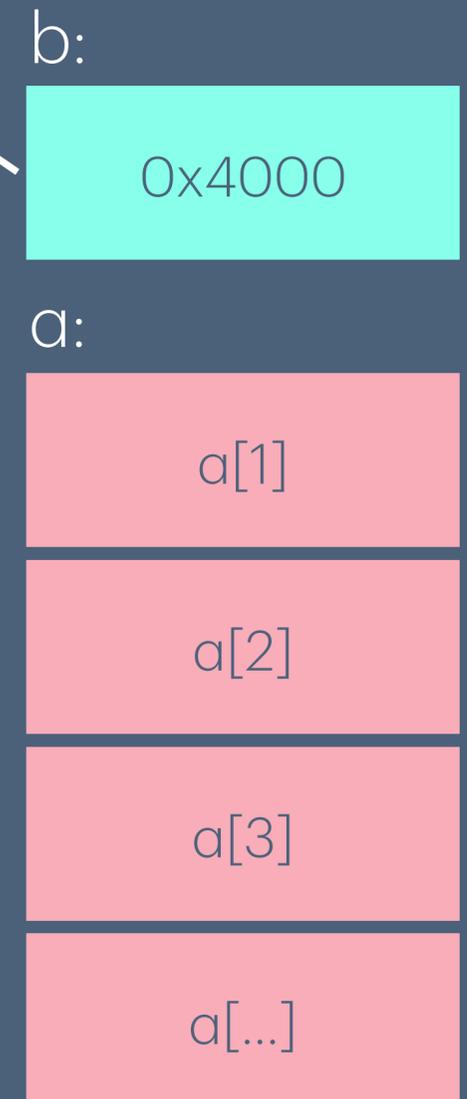
0x4000:



$b + 2 = 0x4008$

$b[2] = *(b + 2) = 0x6767$

0x1000:



Approaches to counting memory reads

Q1: Static & Dynamic Arrays

- Translate C to ASM
 - Pros
 - Reliable
 - Good practice
 - Cons
 - Slower
 - ASM has to be correct
- Analyze the C code
 - Pros
 - Faster
 - Better understanding of C
 - Cons
 - Less flexible

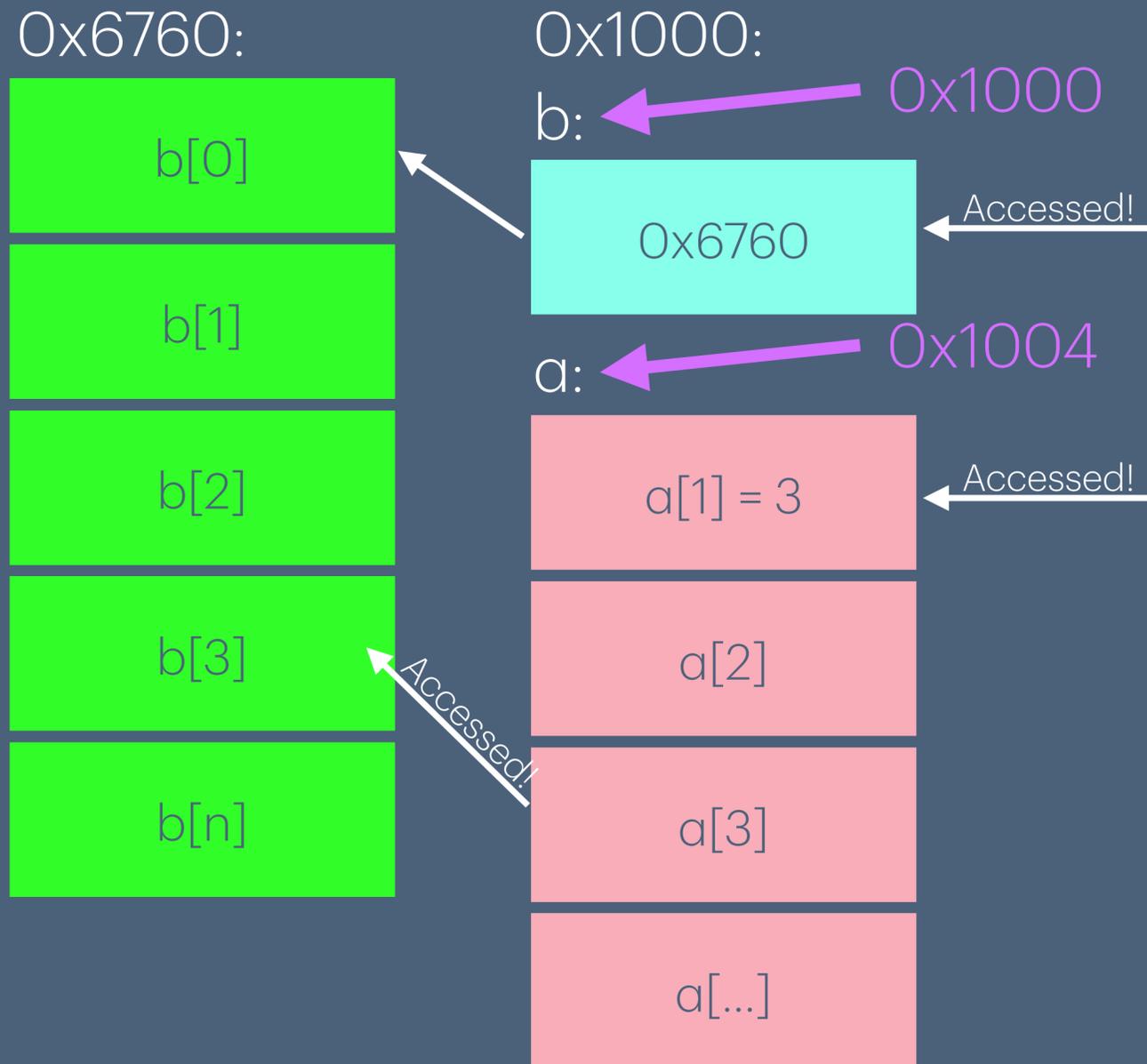
$$b[b[a[1]]] = b[1] + b[a[1]]$$

Dynamic vs Static Arrays

- Analysis approach
 - What do I need?
 - **Address** of LHS
 - **Value** of RHS
 - Values can be shared between RHS and LHS

$$b[b[a[1]]] = b[1] + b[a[1]]$$

Known at compile time

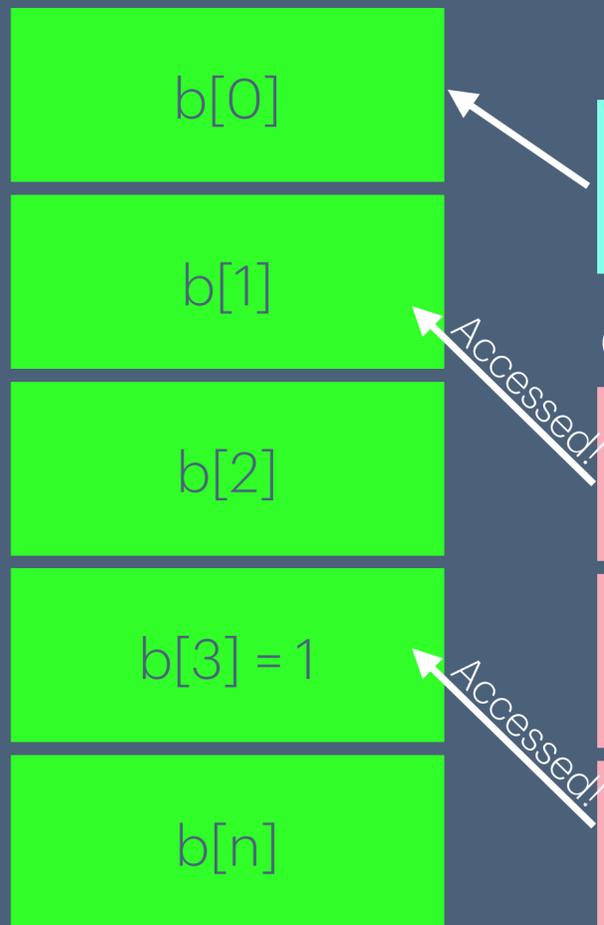


- Address of $b[b[a[1]]]$
- Value of $b[a[1]]$
- Need $a[1]$
 - $a[1]$ known at compile time
 - 1 memory read to get $a[1]$
- Need value of b (where does b point to?)
 - 1 memory read
- Finally, $b[a[1]]$, 1 more memory read

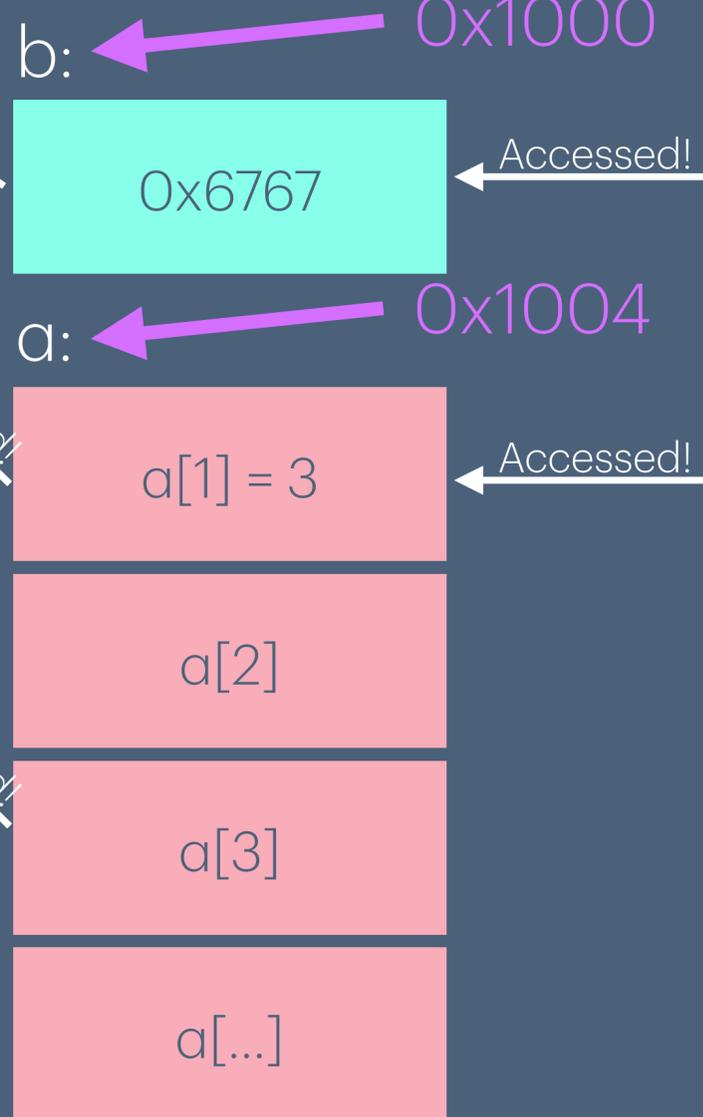
$$b[b[a[1]]] = b[1] + b[a[1]]$$

Known at compile time

0x6760:



0x1000:



- What we know from before

- `a[1]`, `b[a[1]]`, `b`

- Need to find where `b[b[a[1]]]` is still

- Easy, add `b` to `b[a[1]] * 4`, no memory needs needed. $b[b[a[1]]] = 0x6760 + 1 * 4 = 0x6764$.

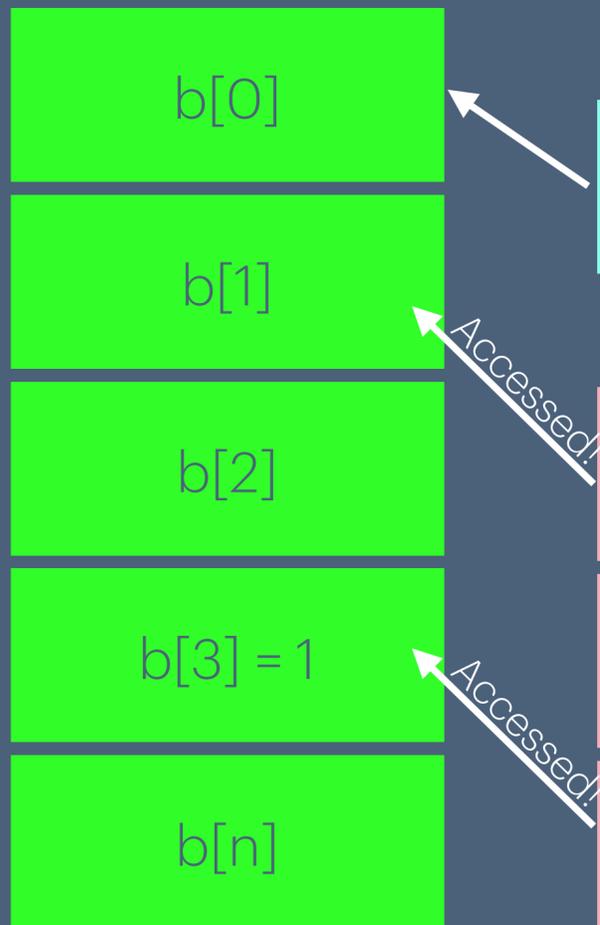
- Calculate `b[1] + b[a[1]]`

- We need `b[1]`, which is 1 memory read

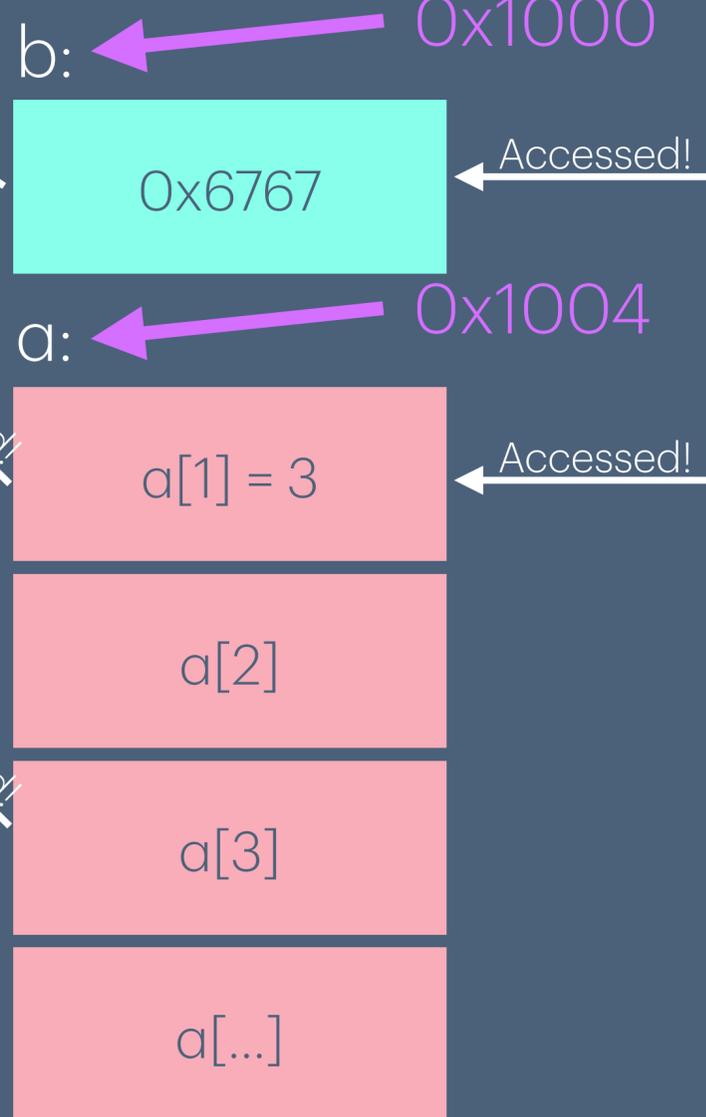
$$b[b[a[1]]] = b[1] + b[a[1]]$$

Known at compile time

0x6760:



0x1000:



- Total memory reads

- 4

- Write assembly approach is probably much slower

- Count number of instructions in ISA that read from memory

Q2: Pointers

WTF?

```
arr[3][0] = arr[0][2 - arr[1][4]]
```

Q2: Pointers

An algorithm

- Apply this rule:

- $\underline{*}(a + \underline{i}) = a[\underline{i}]$

- Work from right to left

- Simplest first

- Let's start with RHS

- `arr[0][2 - arr[1][4]]`

- `arr[0][2 - *(arr[1] + 4)]`

- `arr[0][2 - *(arr[1] + 4)]`

- `arr[0][2 - *(*arr + 1) + 4]`

Q2: Pointers

RHS continued

- `arr[0][2 - *(*arr + 1) + 4]`
 - `*(arr[0] + 2 - *(*arr + 1) + 4)`
- `*(arr[0] + 2 - *(*arr + 1) + 4)`
 - `*(*arr + 0) + 2 - *(*arr + 1) + 4)`
 - `*(*arr + 2 - *(*arr + 1) + 4)`
- DONE!

Q2: Pointers

LHS

- Our original expression: `arr[3][0] = arr[0][2 - arr[1][4]]`
- Convert `arr[3][0]`
 - `arr[3][0]`
 - `*(arr[3] + 0)`
 - `*(arr[3])`
 - `*(*(arr + 3))`
 - `** (arr + 3)`

Q2: Pointers

Grand finale

- `** (arr + 3) = * (*arr + 2 - * (* (arr + 1) + 4))`
- Part 2
 - Find a way to apply this "algorithm" in reverse